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On the weight structure of cyclic codes over GF(q), q > 2.

Abstract

The interrelation between the cyclic structure of an ideal, i.e., a cyclic code over Galois field GF(q), q>2, and its classes of proportional elements is considered. This relation is used in order to define the code's weight structure. The equidistance conditions of irreducible nonprimitive codes over GF(q) are given. Besides that, the minimum distance for some class of nonprimitive cyclic codes is found.

The relation of proportionality for elements of algebra A_n , consisting of polynomials in x over Galois field GF(q), modulo polynomial x^n-1 , is the equivalence relation [1]. Therefore A_n falls into several disjoint subsets and every such subset contains all elements which are proportional to each other. These subsets will be called the classes of proportionality. Let $z(x) \neq 0$ be some vector of A_n . If $\alpha_1 = 1, \alpha_2, \ldots, \alpha_{q-1}$ are all different elements of the multiplicative group $GF(q)^*$ of the field GF(q), then the following q-1 vectors

$$\alpha_1 z(x), \alpha_2 z(x), \dots, \alpha_{q-1} z(x) \tag{1}$$

are some different and proportional to each other elements of A_n . The set of vectors (1) is closed under the multiplication by the elements of the group $GF(q)^*$. Hence the set (1) represents some class of proportional elements, which will be denoted by $P_{z(x)}$. Because of an arbitrary choice for z(x), every nonzero class consists of q-1 elements of the form (1). Consequently A_n contains $(q^n-1)/(q-1)$ different nonzero classes. Evidently all elements of one class have the same period [2, 3] or the same order [8]. Clearly, the supporting sets [2] of vectors, entering into the same class of proportionality, are similar too. Hence, the Hamming weight is also the same for all vectors of one class. Thus, we can say that any proportionality class $P_{z(x)}$, $P_{z(x)} \subset A_n$, has its order, its supporting set and its Hamming weight. Obviously, any proportionality class of A_n is characterized by its unique monic polynomial.

Now consider an ideal J, $J \subset A_n$, i.e., some cyclic code over GF(q), having the following generator $g(x) = (x^n - 1)/h(x)$ [7], where h(x) is some parity-check polynomial of degree m, having some order n, n = ord(h(x)) [8]. Below we suppose that q > 2 and gcd(n, q) = 1.

It's also known [3, 4, 10] that any ideal is partitioned into several disjoint subsets, that is cycles, under the multiplication of ideal's vectors

by x. On the other hand, some ideal J, as a subspace of A_n , consists of $(q^m-1)/(q-1)$ nonzero proportionality classes. Obviously, the existence of two different partitions into some disjoint subsets of any ideal assumes a certain dependence between proportionality classes and cycles of ideal.

Further, any ideal $J \subset A_n$ is the direct sum of minimal ideals [2, 8]

$$J = \sum_{i=1}^{t} J_i, \tag{2}$$

where J_i is some minimal ideal, having an irreducible parity-check polynomial $h_i(x)$ of degree m_i and of order n_i , $n_i = ord(h_i(x))$, $1 \le i \le t$. This implies that the following polynomial

$$h(x) = \prod_{i=1}^{t} h_i(x), \tag{3}$$

is the parity-check polynomial of J and n, $n = lcm(n_1, n_2, ..., n_t)$, is the order of h(x) [8]. It should be stressed that under the condition gcd(n,q) = 1 the polynomial (3) has no repeated factors [8].

Remark 1. It is worth mentioning that the number n can be some number of either the primitive form $n = q^m - 1$ or of the nonprimitive form $n \neq q^m - 1$. In the first case, we have some cyclic primitive code and the second case corresponds to a certain cyclic nonprimitive code [11]. Let us stress that $n \neq q^m - 1$ if and only if $n_i \neq q^{m_i} - 1$, $1 \leq i \leq t$. But if there is at least one primitive polynomial among the polynomials $h_i(x)$, $1 \leq i \leq t$, then $n = q^m - 1$.

Furthermore, applying the theory of linear recurring sequences [8, 12] to elements of an ideal, we obtain that every element of some ideal $J \subset A_n$ is characterized by its unique minimal polynomial. Denote by C the set of all elements of J, having the same minimal polynomial c(x). The set C is either some minimal ideal J_i , $1 \le i \le t$, or a certain subset of all such elements of J, whose characteristic polynomial of the smallest degree coincides with c(x). In the general case, the polynomial c(x) is equal to the product of some k, $1 \le k \le t$, polynomials from t different prime divisors of h(x). Thus,

$$c(x) = \prod_{j=1}^{k} h_{i_j}(x), \quad 1 \le k \le t.$$
 (4)

This means that any element of C has the same period or the same order $n_c = ord(c(x))$, $n_c \le n$, $n_c|n$.

Lemma 1. For the set C, $C \subset J$, having some minimal polynomial c(x) in terms of (4), the following equality takes place

$$n_c \cdot s_c = R_c \cdot (q-1), \tag{5}$$

where s_c is the number of all cycles, and R_c is the number of all proportionality classes of C.

Proof. The set C is closed under two different operations. The first operation is the cyclic shift of vector and the second one is the multiplication of vectors by elements of group $GF(q)^*$. Hence equality (5) can be obtained by the counting of the number of all elements, belonging to C, via the two different ways. The lemma is proved.

Theorem 1. Any cycle $\{z(x)\}$ of ideal (2), having some period n_z , $n_z|n$, consists of d_z subsets. The first element of each such subset is proportional to z(x). Every such subset contains r_z nonproportional to each other vectors, i.e., $n_z = r_z \cdot d_z$, $d_z|(q-1)$. And the number r_z is the index of the subgroup, belonging to $GF(q)^*$, of order d_z in the group of the roots of unity, having the least possible order.

Proof. Let r_z , $1 \le r_z \le n_z$, is the smallest natural number such that the following equality holds

$$x^{r_z} \cdot z(x) = \alpha \cdot z(x) \mod(x^{n_z} - 1), \tag{6}$$

where α is some element of $GF(q)^*$. Then the following r_z vectors of cycle $\{z(x)\}$

$$z(x), x \cdot z(x), \dots, x^{r^z - 1} \cdot z(x) \tag{7}$$

are some non–proportional to each others vectors because, assuming the inverse, we should be able to decrease the number r_z , but it is impossible. Hence the set of elements (7) belongs to the following r_z classes of proportionality

$$P_{z(x)}, P_{xz(x)}, \dots, \dots, P_{x^{r_z-1}z(x)}.$$
 (8)

Since $x^{n_z}z(x)=z(x)$ in the ring A_n and also, considering (6), we see that $P_{z(x)}=P_{x^{r_z}z(x)}=P_{x^{n_z}z(x)}$. This means that the cycle $\{z(x)\}$ belongs to the classes (8) and every such class contains d_z vectors, $d_z=n_z/r_z$, $1< d_z \leq q-1$. In terms of equality (6) the following different vectors $z(x), x^{r_z}z(x), x^{2r_z}z(x), \ldots, x^{(d_z-1)r_z}z(x)$ of class $P_{z(x)}$ can be represented as $\alpha^0 z(x), \alpha z(x), \ldots, \alpha^{d_z-1}z(x)$. This implies that $\alpha^0=1,\alpha,\alpha^2,\ldots,\alpha^{d_z-1}$ are the different elements of group $GF(q)^*$. Since $x^{n_z}z(x)=x^{r_zd_z}z(x)=\alpha^{d_z}z(x)=z(x)$, we see that $\alpha^{d_z}=1$. This yields that d_z is the order of element α in $GF(q)^*$. Consequently, $d_z|q-1$, $1\leq d_z \leq q-1$.

Finally, under the condition gcd(n,q)=1 the polynomial x^n-1 has n different roots in $GF(q^h)$ field, where h is the multiplicative order of q modulo n [8]. Denote by E(n) the multiplicative group of n-th roots of unity over GF(q). Let $\xi \in E(n)$ be some n-th primitive root of unity. Then the following set of elements $\xi^0=1,\xi,\xi^2,\ldots,\xi^{n_z-1},\ldots,\xi^{n-1}$ represents the group E(n). Since $n_z|n$, we have $E(n_z)\subset E(n)$, where $E(n_z)$ is the multiplicative group of n_z -th roots of unity. Moreover,

taking into account the isomorphism of the groups, having the same order [9], we can state that the subgroup $\{\alpha\}$, $\{\alpha\} \subset GF(q)^*$, of order d_z belongs to $E(n_z)$ because $d_z|n_z$.

As mentioned above, n_z is the period of z(x), so that n_z is the smallest divisor of n such that the following congruence $x^{r_z} \equiv \alpha \mod(x^{n_z}-1)$ takes place. Hence $E(n_z)$ is the smallest group of n- th roots of unity, which contains $\{\alpha\}$. Since $\alpha = \xi^{r_z}$, we see that the following elements $\xi^{r_z d_z} = 1, \xi^{r_z}, \xi^{2r_z}, \dots, \xi^{(d_z-1)r_z}$ represent the subgroup $\{\alpha\}$ in the group $E(n_z)$. Besides that, the decomposition of $E(n_z)$ relative to the subgroup $\{\alpha\}$ consists of r_z different cosets. Thus, the number r_z is the index of subgroup $\{\alpha\}$ in the group of roots of unity, having the smallest possible order. The theorem is proved.

Remark 2. Notice that when a parity-check polynomial of code is some primitive polynomial of degree m and of order $n = q^m - 1$, m > 1, then $r = R = (q^m - 1)/(q - 1)$ and d = q - 1. (see [2], [7]).

Corrollary 1. The period n_z , $n_z = r_z \cdot d_z$, of element z(x), $z(x) \in J$, equals d_z , $d_z|(q-1)$, $1 \le d_z \le (q-1)$, if and only if the cycle $\{z(x)\}$ is contained in one class of proportionality.

Corollary 2. All code words of any cyclic code, having some length n over GF(q), fall into some equal-weight subsets and every such subset includes all proportional to each other cycles.

Besides that, consider some minimal ideal J, $J \subset A_n$, having an irreducible nonprimitive parity-check polynomial h(x) of degree m and of order n, $n \neq q^m-1$, i.e., some irreducible code of nonprimitive length.

Remark 3. The degree m of the polynomial h(x) coincides with the multiplicative order h of the number q modulo n [8]. Also, the order n of the polynomial h(x) is some divisor of $q^h - 1$. This means that the order n can change in the following limits $1 < n < q^h - 1$, $n \neq q - 1$. If 1 < n < q - 1, i.e., $n \neq q - 1$, then some minimal ideal J, $J \subset A_n$, of dimention one contains only one nonzero class of proportyonality. Consequently, n = d, 1 < d < q - 1, d|q - 1.

Theorem 2. Any cycle of minimal ideal J, $J \subset A_n$, having some parity-check polynomial h(x) of degree m, m > 1, and of order n, $n \neq q^m - 1$, is contained in r proportionality classes, $1 \leq r \leq R$, r|R, $R = (q^m - 1)/(q - 1)$. Every such class consists of d, $1 \leq d \leq (q - 1)$, different vectors of cycle, and

$$d = qcd(q - 1, n), \quad r = n/d. \tag{9}$$

Proof. All elements of minimal ideal J, $J \subset A_n$, have the same order n, n = ord(h(x)). Applying the theorem 1 to some element f(x), $f(x) \in J$, we have $n = r_f \cdot d_f$, $d_f|_{q-1}$, $1 \le r_f \le n$, $1 \le d_f \le (q-1)$, and also the following equality

$$x^{r_f} \cdot f(x) = \gamma \cdot f(x), \tag{10}$$

where γ is some element of $GF(q)^*$. Evidently, if either $r_f = 1$, i.e., $d_f = n$, or $r_f = n$ and $d_f = 1$, then equalities (9) take place. Hence, below we suppose that $1 < r_f < n$, $1 < d_f < q - 1$, and therefore $q - 1 < n < q^m - 1$.

Taking into account (9), we have $d_f \leq d$. Let us show that the strong unequality $d_f < d$ is impossible. Indeed, if $d_f < d$, then the subgroup $\{\gamma\}$, where γ is the element from equality (10), belongs to some group of n-th roots of unity, having the order d, because $d_f|d$. Since d|(q-1), we see that the subgroup $\{\gamma\}$ belongs to $GF(q)^*$. This means that the cycle $\{f(x)\}$ is contained in one class of proportionality, i. e., $r_f = 1$. But this fact contradicts to the condition $r_f > 1$. This implies that the strong unequality $d_f < d$ is impossible. Hence $d_f = d$. Because of an arbitrary choice of f(x) we can conclude that equalities (9) take place for any element of J. The theorem is proved [6].

Remark 4. It is necessary to note that the theorem 2 is valid only for irreducible codes of non–primitive length except Reed-Solomon codes of length n=q-1 as it was shown above. In the case of irreducible codes of primitive length n, $n=q^m-1$, m>1, the theorem 2 will be valid if and only if gcd(m,q-1)=1. Indeed, when the last condition takes place, then $gcd((q^m-1)/(q-1),q-1)=1$ [7]. Thus, $d=q-1=gcd(q^m-1,q-1)$ that is d=gcd(n,q-1). It follows that the theorem 2 holds.

Remark 5. Notice that under the condition gcd(m, q-1) = 1 the number r from (9) has no divisors of (q-1) except 1, so gcd(r, q-1) = 1. This means that gcd(r, d) = 1. Hence, considering the fact that r|R and also, taking into account (9) and the following equality s = R(q-1)/n, we have r = gcd(R, n). Besides that, if either one from the two numbers n and q-1 does not contain multiple prime divisors or the same prime divisors of these numbers have the same degrees under the decomposition of both n and q-1, then the following equalities r = gcd(R, n), gcd(r, d) = 1 also take place.

Corollary 3. Both the number r and the number d are the same numbers of all irreducible divisors of polynomial $x^n - 1$ over GF(q), having the same order.

Corollary 4. The number R, $R = (q^m - 1)/(q - 1)$, of proportionality classes, of some irreducible code K, having some length n, $n \neq (q^m - 1)$, $n = d \cdot r$, over GF(q) field, consists of some v different subsets. And every such subset contains r equal-weight proportionality classes, i.e., $R = v \cdot r$. Besides that, every subset includes b equal-weight cycles, b = (q - 1)/d, $1 < b \le (q - 1)$. So that the number of all cycles for K equals $s = v \cdot b$ and qcd(r, b) = 1.

Proof. According to the theorem 2 any cycle of code K is contained in r, r|R, proportionality classes. Therefore the number v=R/r gives us the common quantity of different subsets of J, each of which consists of r classes, i.e., $R=v\cdot r$. The number b, b=(q-1)/d, is the number

of all different equal-weight cycles, contained in every such subset, which consists of some r classes. Hence the number of all cycles for K is equal to $s = v \cdot b$. Since $d = \gcd(q-1,n)$ we see that $\gcd(r,b) = 1$. Actually, assuming the inverse, we would have been able to decrease the number d, but it's impossible. The corollary is proved.

Corollary 5. The irreducible nonprimitive code K is some equidistant code if s = b. Besides that, the last equation is equivalent to the following ones: r = R or gcd(s, R) = 1.

Remark 6. Note that the condition s = b was obtained in [14, 15], but only for some subclass of irreducible nonprimitive codes and under the following additional restriction gcd(b, m) = 1.

Corollary 6. The weight of any element, belonging to some irreducible nonprimitive code K of length n over GF(q), is multiple of the number d, $d = \gcd(q-1,n)$.

Proof. The weight of any element z(x), $z(x) \in K$, of order n, n = ord(h(x)), is equal to the number of such j, $0 \le j \le n-1$, for which the polynomial $x^j \cdot z(x)$ has the following degree n-1. According to the theorem 2, the number of such polynomials for the cycle z(x), having degree n-1, is equal to $w_r \cdot d$, where w_r is the number of polynomials, having degree n-1, among the first r cyclic shifts of z(x), and d = gcd(q-1,n). The corollary is proved.

In addition, consider some ideal J, $J \subset A_n$, of the form (2), having the parity-check polynomial h(x) in terms of (3).

Theorem 3.If the following condition gcd(h, q-1) = 1, where h is the multiplicative order of number $q \mod n$, takes place, then any cycle of set C, $C \subset J$, having some minimal polynomial of the form (4), is contained in r_c proportionality classes, $r_c|R_c$, and every such class includes d_c , $d_c|q-1$, elements of cycle, that is $n_c = r_c \cdot d_c$, $n_c = ord(c(x))$, where R_c is the number of all proportionality classes of set C, and

$$r_c = \gcd(R_c, n_c), \quad d_c = \gcd(q - 1, n_c). \tag{11}$$

Proof. It is sufficient to consider the case k=2 because the general case can be obtained by the induction. Thus assume that $c(x)=h_1(x)\cdot h_2(x)$, where $h_i(x)$ is of degree m_i and of order n_i , $n_i=q^{m_i}-1$, $1\leq i\leq 2$, is a certain prime miltiplier of c(x). It is known [8], that the number m_i , $1\leq i\leq 2$, equals either h or some divisor of this number. Hence, taking into account the theorem 2, and also remarks 4 and 5, we have $n_i=r_i\cdot d_i$, where $r_i=\gcd(R_i,n_i)$, $d_i=\gcd(q-1,n_i)$, $\gcd(r_i,q-1)=1$, $1\leq r_i\leq R_i$, $1\leq d_i\leq q-1$, and $R_i=(q^{m_i}-1)/(q-1)$, where R_i is the number of all proportionality classes of minimal ideal J_i , $1\leq i\leq 2$. Therefore the order n_c , $n_c=lcm(n_1,n_2)=n_1\cdot n_2/\gcd(n_1,n_2)$ of the polynomial c(x) can be rewritten as

$$n_c = r_1 \ d_1 \cdot r_2 \ d_2/gcd(r_1d_1 \cdot r_2d_2). \tag{12}$$

Since $gcd(r_i, q - 1) = 1$, we have $gcd(r_i, d_1 \cdot d_2) = 1$, $1 \le i \le 2$. Thus $gcd(r_1r_2, d_1d_2) = 1$. Hence $gcd(gcd(r_1, r_2), gcd(d_1, d_2)) = 1$ so that (12) may be represented in the following form

$$n_c = (r_1 \ r_2/gcd(r_1, r_2)) \cdot (d_1 \ d_2/gcd(d_1, d_2)). \tag{13}$$

Thus, $n_c = lcm(r_1, r_2) \cdot lcm(d_1, d_2)$. Now by r_c and d_c we denote $lcm(r_1, r_2)$ and $lcm(d_1, d_2)$, respectively. Thus $n_c = r_c \cdot d_c$ and $gcd(r_c, d_c) = 1$. Since $d_c|q-1$ and $gcd(r_c, q-1) = 1$, we obtain $d_c = gcd(q-1, n_c)$. Considering (5), it follows that $n_c|(R_c \cdot (q-1))$. Hence we have $r_c = gcd(R_c, n_c)$. Consequently the order n_c of any element of set C is equal to the product of two relatively prime numbers, i.e., $n_c = r_c \cdot d_c$, where $r_c = lcm(r_1, r_2) = gcd(R_c, n_c)$ and $d_c = lcm(d_1, d_2) = gcd(q-1, n_c)$.

Furthermore, applying the theorem 1 to some element $a(x) \in C$ of period $n_c = r_a \cdot d_a$, we have

$$x^{r_a} \cdot a(x) = \theta a(x), \tag{14}$$

where $\theta \in GF(q)^*$ is some element of order d_a , and r_a is the smallest natural number such that equality (14) takes place. Notice that the subgroup $\{\theta\}$ has the order d_a , $d_a < d_c$. If $d_c = 1$, then $d_a = 1$ and $n_c = r_a = r_c = \gcd(R_c, n_c)$, so that equalities (11) hold. For this reason below we suppose that $d_c > 1$. If under this condition the number r_c is equal to one, then $n_c = d_c = \gcd(q - 1, n_c)$ and the theorem is valid. Therefore below we suppose that both $r_c > 1$ and $d_c > 1$.

Evidently, $d_a \leq d_c$. Now let us show that the inequality $d_a < d_c$ is not possible. Indeed, if $d_a < d_c$, then we come to the following conclusion. The subgroup $\{\theta\}$, where θ is the element from equality (14), belongs to some subgroup of $GF(q)^*$, having the order d_c , because $d_a|d_c$. Since d_c is some divisor of n_c , then, considering the uniqueness of subgroups, having the same order, the subgroup $\{\theta\}$ belongs to some group of d_c -th roots of unity. This implies that the smallest group of roots of unit, containing $\{\theta\}$, has an order, which either less or equals d_c . Thus, both the period of a(x) and the order of c(x) must be either less or equal to d_c . This yields that the order of c(x) must be some divisor of q-1. But this fact contradicts the condition $r_c > 1$. Hence the inequality $d_a < d_c$ is impossible so that $d_a = d_c$ and $r_a = r_c$. Because of an arbitrary choice of a(x) equalities (11) take place for any element of set C. The theorem is proved.

Corollary 7. The order n_c of reducible factor c(x) of polynomial $x^n - 1$, having some degree m over GF(q), is some divisor of number $q^m - 1$, if gcd(h, q - 1) = 1, where h is the multiplicative order of number $q \mod n$.

Remark 7. In terms of condition gcd(h, q - 1) = 1, where h is the multiplicative order of number $q \mod n$, the theorem 3 is valid for cyclic codes of both the primitive and the nonprimitive length. Also,

taking into consideration the remark 5, the order of any reducible factor of the polynomial $(x^n - 1)$ over GF(q) of degree m, is some divisor of the number $(q^m - 1)$, if gcd(h, q - 1) = 1.

Theorem 4. Any cycle of set C, $C \subset J$, having some minimal polynomial c(x) of the type (4) and of order $n_c = lcm(n_1, n_2, ..., n_k)$, where $n_i \neq q^{m_i} - 1$, $1 \leq i \leq k$, is contained in r_c , $r_c | R_c$, $1 \leq r_c \leq R_c$, classes of proportionality and every such class includes d_c , $1 \leq d_c \leq q-1$, elements of cycle, where R_c is the number of all proportionality classes of C, and

$$d_c = gcd(q - 1, n_c), \quad r_c = n_c/d_c.$$
 (15)

Proof. It is sufficient to assume that k=2 because the general case can be obtained by the induction. This implies that $c(x) = h_1(x)h_2(x)$, where $h_i(x)$ is some prime divisor of equality (3), having some degree m_i , and of order n_i , $n_i \neq q^{m_i} - 1$, $1 \leq i \leq 2$. This yields that the theorem 2 holds for the minimal ideal J_i , $1 \leq i \leq 2$.

Evidently, if one of the two numbers, i.e., either d_c or r_c is equal to one, then equalities (15) hold. For this reason below we assume that both $r_c > 1$ and $d_c > 1$.

Let z(x) be some vector of set C. According to the theorem 1 some cycle $\{z(x)\}\subset C$ is contained in r_z classes and every such class includes d_z different elements of this cycle. Thus $n_c=r_z$ d_z , $1\leq d_z\leq q-1$, $d_z|q-1$, $1\leq r_z\leq n_c$. And in addition, the following equality takes place

$$x^{r_z} \cdot z(x) = \beta \ z(x), \tag{16}$$

where β is some element of $GF(q)^*$ and the subgroup $\{\beta\}$, $\{\beta\} \subset GF(q)^*$, has the order d_z . Evidently, $d_z \leq d_c$. Let us show that the number d_z can not be smaller than d_c . Assume the inverse, i.e. let d_z be less than d_c . Since at least one of the two numbers ether r_1 or r_2 is not equal to one, we see that at least one of the numbers n_i , $1 \leq i \leq 2$, is more than q-1, as was established in the theorem 2. This means that

$$n_c = lcm(n_1, n_2) > (q - 1).$$
 (17)

Further, since $d_z|d_c$, we see that the subgroup $\{\beta\}$ of order d_z belongs to some subgroup of $GF(q)^*$, having the order d_c , where β is the element from equality (16). Due to the uniqueness of groups, having the same order, the subgroup $\{\beta\}$ belongs to the group of d_c -th roots of unity because $d_c|n_c$. It follows that the smallest group of the roots of unity, which contains the subgroup $\{\beta\}$, has the order less or equals to d_c . This implies that both the period of z(x) and the order of c(x) is some divisor of q-1. But this fact contradicts to (17). It follows that our assumption is not true, so that $d_z=d_c$, and $r_z=r_c$.

Besides that, since the number d_c is the same number for every cycle of set C, we see that every subset, consisting of r_c proportionality classes, contains the same number of cycles, which is equal to $b_c = q - 1/d_c$.

Moreover, since $d_c = \gcd(q-1, n_c)$, we obtain $\gcd(b_c, r_c) = 1$. Hence, taking into account the following equality $s_c = R_c(q-1)/n_c$, which follows from (5), we see that r_c is some divisor of R_c and $s_c = v_c b_c$, where v_c is equal to R_c/r_c . The theorem is proved.

Corollary 7. (Equidistance signs of the subset C, $C \subset J$, having some minimal polynomial of the form (4).)

All vectors of the subset C, $C \subset J$, having some minimal polynomial c(x) of order $n_c = lcm(n_1, n_2, ..., n_k)$, $n_i \neq q^{m_i} - 1$, $1 \leq i \leq k$, k > 1, have the same weight if at least one of the following conditions holds: 1. $s_c = b_c$, $1 < b_c \leq q - 1$, 2. $r_c = R_c$, 3. $gcd(s_c, R_c) = 1$.

Corollary 8. The order of the reducible factor c(x) of the polynomial x^n-1 over GF(q), having some degree m, is some divisor of the number q^m-1 , if the decomposition of the polynomial c(x) into prime multiples does not contain any primitive polynomial.

Remark 8. Note that corollaries (6) and (8) show us in what cases corollary 3.4 from [8] takes place for some reducible polynomial of the degree m over GF(q).

Also, consider some cyclic code, having the following parity-check polynomial

$$h(x) = \prod_{i=1}^{t} h_i(x),$$
 (18)

where $h_i(x)$ is an irreducible polynomial over GF(q), q > 2, of degree m_i and of order $n_i = (q^{m_i} - 1)/b_i$, $b_i|q - 1$, $1 \le b_i \le q - 1$, $1 \le i \le t$, $gcd(m_i, m_j) = 1$ and $gcd(b_i, b_j) = 1$, provided $i \ne j$, $1 \le i, j \le t$, so that the order n of h(x) equals $n = lcm(n_1, n_2, ..., n_t)$.

It is worth mentioning that in [14] and [15] the following cases of polynomial (18) have been considered. Namely t = 1 and t = 2, and besides that, with some additional restrictions, which can be omitted. Also, some particular case of polynomial (18), that is provided $b_i = 1$ for all i, $1 \le i \le t$, was obtained in [16]. But there are some unnecessary restrictions in this paper too. Also, there are some mistakes in that paper. Namely, the order of the product for some two polynomials from (18) was defined incorrectly in [16].

Finally, using the results obtained above, we have found the minimal distance of code, having the parity-check polynomial (18), (see [17]).

Denote by M the set of degrees for the polynomials $h_i(x)$, $1 \le i \le t$, from the eq. (18) that is $M = (m_1, m_2, ..., m_t)$. Let the number m denote the degree of polynomial h(x), i. e., $m = m_1 + m_2 + ... + m_t$. By $M_{j,k}$, $1 \le j \le C_t^k$, $1 \le k \le t - 1$, denote j's k-subset of M, where C_t^k is the binomial coefficient.

Thus $M_{j,k}=(m_{j_1},m_{j_2},...,m_{j_k})$, $1\leq j\leq C_t^k$, $1\leq k\leq t-1$. At last, denote by $m_{j,k}$ the sum of degrees, from the subset $M_{j,k}$, so that $m_{j,k}=m_{j_1}+m_{j_2}+...+m_{j_k}$, $1\leq j\leq C_t^k$, $1\leq k\leq t-1$. It is obvious that $1\leq m_{j_k}< m$. Let us remark that provided k=1 the number

 $m_{j,k} = m_{j,1} = m_j$ and $1 \le j \le t$ because $C_t^1 = t$.

Theorem 5. The minimal distance of code, having the parity-check polynomial (18), has the following form

$$d_{min} = q^{m-1} - \sum_{j=1}^{C_t^1} q^{m_{j,1}-1} - \sum_{j=1}^{C_t^2} q^{m_{j,2}-1} - \dots - \sum_{j=1}^{C_t^{t-1}} q^{m_{j,t-1}-1}, t \ge 2,$$

$$d_{min} = q^{m-1}(q-1)/b, t = 1, 1 \le b \le q-1,$$

$$d_{min} = q^{m-1}(q-1), t = 1, b = 1.$$

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